

# Convex Duality and Interior Point Method

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# Outline

- Convexity
- Duality
- Optimality
- Examples: Linear Prog. and SVM
- (Interior Point Method)

# Convex Set

A set  $S \in R^n$  is said convex if  
 $\forall x, y \in S$ , and  $0 \leq \theta \leq 1$ ,

$$\theta x + (1 - \theta)y \in S$$

Examples:

- $\{x \in R^n : |x - x_0| < r, \text{ for some } x_0 \in R^n, \text{ and some } r > 0\}$
- $\{x \in R^n : a_i^T x \geq b, \quad i = 1, \dots, m\}$

# Thm 1: $\bigcap_{\alpha} S_{\alpha}$ is convex

## Theorem 1

If  $\{S_{\alpha}\}$  is convex, then  $\bigcap_{\alpha} S_{\alpha}$  is convex.

proof:

$$x, y \in \bigcap_{\alpha} S_{\alpha} \Rightarrow x, y \in S_{\alpha}, \forall \alpha.$$

$$S_{\alpha} \text{ is convex} \Rightarrow \forall \theta \in [0, 1], \theta x + (1 - \theta)y \in S_{\alpha}, \forall \alpha$$

$$\Rightarrow \theta x + (1 - \theta)y \in \bigcap_{\alpha} S_{\alpha}$$

# Convex Function

A function  $f : R^n \rightarrow R$  is said convex if

1.  $\text{dom } f$  is convex
2.  $\forall x, y \in \text{dom } f$ , and  $0 \leq \theta \leq 1$ ,

$$f(\theta x + (1 - \theta)y) \leq \theta f(x) + (1 - \theta)f(y)$$

remark:  $f$  is concave if  $-f$  is convex.

# Optimization Problem

General optimization is in the form: (primal problem)

$$\min f_0(x)$$

$$\text{s.t. } f_i(x) \leq 0, i = 1, \dots, m$$

$$h_i(x) = 0, i = 1, \dots, p$$

We make NO assumption about the convexity.

# Lagrangian $L(x, \lambda, \nu)$

Lagrangian is defined as

$L : R^n \times R^m \times R^p \rightarrow R$ :

$$L(x, \lambda, \nu) = f_0(x) + \sum_{i=1}^m \lambda_i f_i(x) + \sum_{i=1}^p \nu_i h_i(x)$$

We call  $\lambda, \nu$  dual variables, or Lagrange multiplier vector.

# Lagrange Dual Function $g(\lambda, \nu)$

Lagrangian dual function is defined as

$g : R^m \times R^p \rightarrow R$ :

$$g(\lambda, \nu) = \inf_{x \in D} \left\{ f_0(x) + \sum_{i=1}^m \lambda_i f_i(x) + \sum_{i=1}^p \nu_i h_i(x) \right\}$$

## Thm 2: $g(\lambda, \nu)$ is concave

### Theorem 2

$g(\lambda, \nu)$  is concave.

Proof:

First, claim that  $\inf(f + g) \geq \inf f + \inf g$ .

proof: Use the fact

1.  $\sup f = -\inf(-f)$ .

2.  $\sup(f + g) \leq \sup f + \sup g$

## $g(\lambda, \nu)$ is concave (2)

Next, consider  $\phi_i = (\lambda_i, \nu_i) \in \text{dom}(-g)$ ,  $i = 1, 2$ .  
We can show that

$$\begin{aligned} -g(\theta\phi_1 + (1 - \theta)\phi_2) &= \dots \\ &\leq -\{\theta g(\phi_1) + (1 - \theta)g(\phi_2)\} \end{aligned}$$

Thus  $-g(\lambda, \nu)$  is convex, hence  $g(\lambda, \nu)$  is concave.

## Thm 3: $g(\lambda, \nu) \leq p^*$

### Theorem 3

If  $p^*$  is the optimal value of the primal problem, then  $\forall \lambda \geq 0$ , and any  $\nu$ , we have

$$g(\lambda, \nu) \leq p^*.$$

**Proof:** Since  $\lambda_i \geq 0$ ,  $f_i(x^*) \leq 0$ ,  $h_i(x^*) = 0$ , we have  $\sum_{i=1}^m \lambda_i f_i(x^*) + \sum_{i=1}^p \nu_i h_i(x^*) \leq 0$ . So,

$$\begin{aligned} g(\lambda, \nu) &= \inf L(x, \lambda, \nu) \leq L(x^*, \lambda, \nu) \\ &= f_0(x^*) + \sum_{i=1}^m \lambda_i f_i(x^*) + \sum_{i=1}^p \nu_i h_i(x^*) \leq f_0(x^*) \end{aligned}$$

# Dual Problem

Theorem 3 says  $g(\lambda, \nu)$  is a lower bound.  
 $\Rightarrow$  want to find the best lower bound!

Dual problem:

$$\begin{array}{ll} \max_{\lambda, \nu} & g(\lambda, \nu) \\ \text{s.t.} & \lambda \geq 0 \end{array}$$

This is convex!

# Strong and Weak Duality

Let  $p^*$  be the optimal value of primal, and  $d^*$  of dual, then since  $g(\lambda, \nu) \leq p^*, \forall \lambda, \nu$  (Thm 3)

**Weak duality:**

$$d^* \leq p^*$$

always hold.

If

$$d^* = p^*$$

then we say **Strong duality** holds (e.g. when primal constraints are convex.)

# Geometric Interpretation

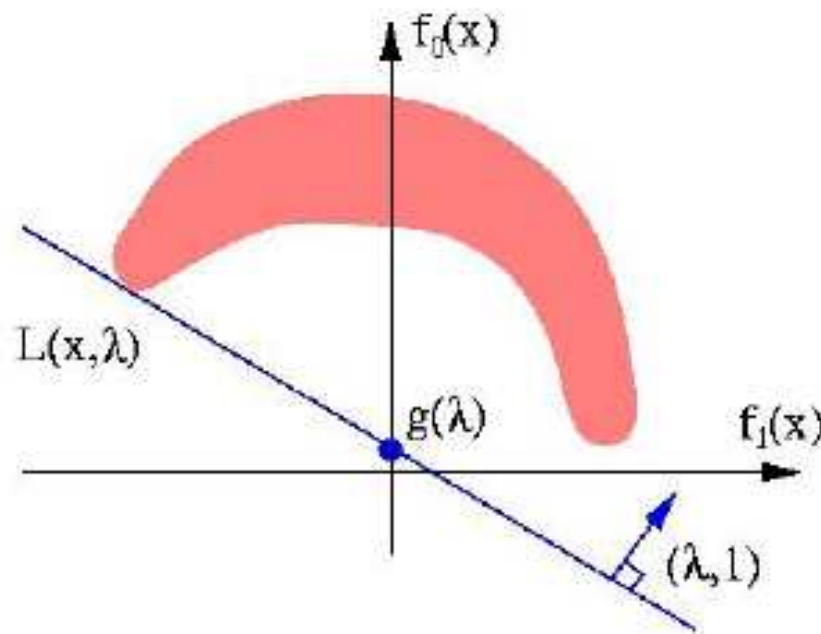


Figure 1: Blue lines is  $L(x, \lambda) = f_0(x) + \lambda f_1(x)$  and  $\lambda$  is the slope.  $g(\lambda) = \inf L$  is achieved at  $f_1(x) = 0$ .

# Geometric Interpretation (2)

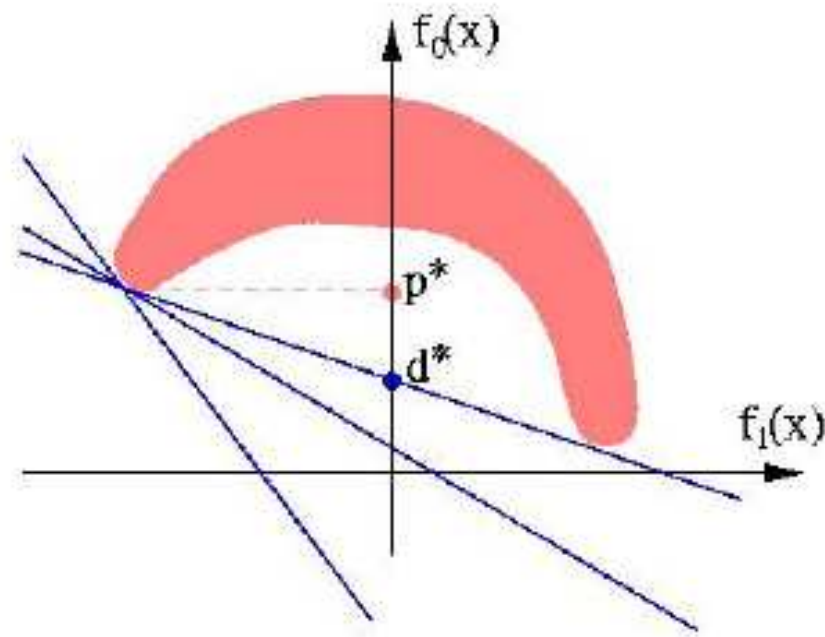


Figure 2: If primal is not convex, then  $p^* \neq d^*$

# KKT Condition

KKT: If  $x^*$  is a minimizer, then

1. (feasibility)  $f_i(x^*) \leq 0, h_i(x^*) \leq 0$

2. (dual feasible)  $\lambda_i^* \geq 0$

3. (complementary slackness)  $\lambda_i^* f_i(x^*) = 0$

4. (first order stationarity)

$$\nabla f_0(x^*) + \sum \lambda_i \nabla f_i(x^*) + \sum \nu_i \nabla h_i(x^*) = 0.$$

The converse is true only if the primal is convex.

# Why dual?

1. Dual optimal = primal optimal if primal is convex
2. Dual problem is always convex
3. Convex problems can be solved efficiently via interior method

# Example 1: Linear Program

Linear Programming in standard form:

$$\begin{aligned} \min \quad & c^T x \\ \text{s.t.} \quad & Ax = b \\ & x \geq 0 \end{aligned}$$

Let's start from the Lagrangian:

$$\begin{aligned} L(x, \lambda, \nu) &= c^T x - \lambda^T x + \nu^T (Ax - b) \\ &= -b^T \nu + (c + A^T \nu - \lambda)^T x \end{aligned}$$

## Example 1: Linear Prog. (2)

Now, having the Lagrangian,

$$L(x, \lambda, \nu) = -b^T \nu + (c + A^T \nu - \lambda)^T x,$$

we can compute Lagrange dual function

$$\begin{aligned} g(\lambda, \nu) &= \inf_{x \in D} L(x, \lambda, \nu) \\ &= -b^T \nu + \inf_{x \in D} \{(c + A^T \nu - \lambda)^T x\} \\ &= \begin{cases} -b^T \nu & : \quad c + A^T \nu - \lambda = 0, \\ -\infty & : \quad \text{otherwise.} \end{cases} \end{aligned}$$

# Example 1: Linear Prog. (3)

Therefore, the dual problem is:

$$\begin{aligned} \max_{\nu, \lambda} \quad & -b^T \nu \\ \text{s.t.} \quad & c + A^T \nu - \lambda = 0, \quad \lambda \geq 0 \end{aligned}$$

Equivalent to:

$$\begin{aligned} \max_{\nu} \quad & -b^T \nu \\ \text{s.t.} \quad & c + A^T \nu \geq 0 \end{aligned}$$

# Example 2: SV classification

Optimal Hyperplane problem:

$$\begin{aligned} \min_{\mathbf{w}, b} \quad & \frac{1}{2} \|\mathbf{w}\|^2 \\ \text{s.t.} \quad & y_i(\mathbf{w}^T \mathbf{x}_i + b) \geq 1, \quad i = 1, \dots, m. \end{aligned}$$

The Lagrangian is

$$L((\mathbf{w}, b), \lambda) = \frac{1}{2} \|\mathbf{w}\|^2 + \sum_{i=1}^m \lambda_i (1 - y_i(\mathbf{w}^T \mathbf{x}_i + b))$$

## Example 2: SV classification (2)

Lagrange dual function is

$$g(\lambda) = \inf_{\mathbf{w}, b} L((\mathbf{w}, b), \lambda)$$

The infimum can be achieved when

$$\frac{\partial}{\partial \mathbf{w}} L = 0 \quad \text{and} \quad \frac{\partial}{\partial b} L = 0.$$

This leads to

$$\mathbf{w} = \sum_{i=1}^m \lambda_i y_i \mathbf{x}_i \quad \text{and} \quad \sum_{i=1}^m \lambda_i y_i = 0$$

## Example 2: SV classification (3)

So we have

$$g(\lambda) = \sum_{i=1}^m \lambda_i - \frac{1}{2} \sum_{i,j=1}^m \lambda_i \lambda_j y_i y_j \mathbf{x}_i^T \mathbf{x}_j$$

Hence the dual problem is:

$$\max_{\lambda} \quad \sum_{i=1}^m \lambda_i - \frac{1}{2} \sum_{i,j=1}^m \lambda_i \lambda_j y_i y_j \mathbf{x}_i^T \mathbf{x}_j$$

$$\text{s.t.} \quad \sum_{i=1}^m \lambda_i y_i = 0, \quad \lambda \geq 0.$$

# Questions?

Thank you.